Research Article

e-Duality Theorems for Convex Semidefinite Optimization Problems with Conic Constraints

Gue Myung Lee and Jae Hyoung Lee

Department of Applied Mathematics, Pukyong National University, Pusan 608-737, South Korea

Correspondence should be addressed to Gue Myung Lee, gmlee@pknu.ac.kr

Received 30 October 2009; Accepted 10 December 2009

Academic Editor: Yeol Je Cho

Copyright © 2010 G. M. Lee and J. H. Lee. This is an open access article distributed under the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original work is properly cited.

A convex semidefinite optimization problem with a conic constraint is considered. We formulate a Wolfe-type dual problem for the problem for its ϵ -approximate solutions, and then we prove ϵ -weak duality theorem and ϵ -strong duality theorem which hold between the problem and its Wolfe type dual problem. Moreover, we give an example illustrating the duality theorems.

1. Introduction

Convex semidefinite optimization problem is to optimize an objective convex function over a linear matrix inequality. When the objective function is linear and the corresponding matrices are diagonal, this problem becomes a linear optimization problem.

For convex semidefinite optimization problem, Lagrangean duality without constraint qualification [1, 2], complete dual characterization conditions of solutions [1, 3, 4], saddle point theorems [5], and characterizations of optimal solution sets [6, 7] have been investigated.

To get the *e*-approximate solution, many authors have established *e*-optimality conditions, *e*-saddle point theorems and *e*-duality theorems for several kinds of optimization problems [1, 8-16].

Recently, Jeyakumar and Glover [11] gave *e*-optimality conditions for convex optimization problems, which hold without any constraint qualification. Yokoyama and Shiraishi [16] gave a special case of convex optimization problem which satisfies *e*-optimality conditions. Kim and Lee [12] proved sequential *e*-saddle point theorems and *e*-duality theorems for convex semidefinite optimization problems which have not conic constraints.

The purpose of this paper is to extend the e-duality theorems by Kim and Lee [12] to convex semidefinite optimization problems with conic constraints. We formulate a Wolfe type dual problem for the problem for its e-approximate solutions, and then prove

e-weak duality theorem and *e*-strong duality theorem for the problem and its Wolfe type dual problem, which hold under a weakened constraint qualification. Moreover, we give an example illustrating the duality theorems.

2. Preliminaries

Consider the following convex semidefinite optimization problem:

(SDP) Minimize
$$f(x)$$
,
subject to $F_0 + \sum_{i=1}^m x_i F_i \ge 0$, $(x_1, x_2, \dots, x_m) \in C$, (2.1)

where $f : \mathbb{R}^m \to \mathbb{R}$ is a convex function, *C* is a closed convex cone of \mathbb{R}^m , and for $i = 0, 1, ..., m, F_i \in S_n$, where S_n is the space of $n \times n$ real symmetric matrices. The space S_n is partially ordered by the Löwner order, that is, for $M, N \in S_n, M \geq N$ if and only if M - N is positive semidefinite. The inner product in S_n is defined by (M, N) = Tr[MN], where $\text{Tr}[\cdot]$ is the trace operation.

Let $S := \{M \in S_n \mid M \succeq 0\}$. Then *S* is self-dual, that is,

$$S^{+} := \{ \theta \in S_{n} \mid (\theta, Z) \ge 0, \text{ for any } Z \in S \} = S.$$

$$(2.2)$$

Let $F(x) := F_0 + \sum_{i=1}^m x_i F_i$, $\hat{F}(x) := \sum_{i=1}^m x_i F_i$, $x = (x_1, \dots, x_m) \in \mathbb{R}^m$. Then \hat{F} is a linear operator from \mathbb{R}^m to S_n and its dual is defined by

$$\widehat{F}^*(Z) = (\operatorname{Tr}[F_1 Z], \dots, \operatorname{Tr}[F_m Z]),$$
(2.3)

for any $Z \in S_n$. Clearly, $A := \{x \in C \mid F(x) \in S\}$ is the feasible set of SDP.

Definition 2.1. Let $g : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$ be a convex function.

(1) The subdifferential of *g* at $a \in \text{dom } g$, where dom $g = \{x \in \mathbb{R}^n \mid g(x) < +\infty\}$, is given by

$$\partial g(a) = \{ v \in \mathbb{R}^n \mid g(x) \ge g(a) + \langle v, x - a \rangle, \ \forall x \in \mathbb{R}^n \},$$
(2.4)

where $\langle \cdot, \cdot \rangle$ is the scalar product on \mathbb{R}^n .

(2) The ϵ -subdifferential of g at $a \in \text{dom } g$ is given by

$$\partial_{\varepsilon}g(a) = \{ v \in \mathbb{R}^n \mid g(x) \ge g(a) + \langle v, x - a \rangle - \varepsilon, \ \forall x \in \mathbb{R}^n \}.$$

$$(2.5)$$

Definition 2.2. Let $e \ge 0$. Then $\overline{x} \in A$ is called an *e*-approximate solution of SDP, if, for any $x \in A$,

$$f(x) \ge f(\overline{x}) - \epsilon. \tag{2.6}$$

Definition 2.3. The conjugate function of a function $g : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$ is defined by

$$g^*(v) = \sup\{\langle v, x \rangle - g(x) \mid x \in \mathbb{R}^n\}.$$
(2.7)

Definition 2.4. The epigraph of a function $g : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$, *epi* g, is defined by

$$epig = \{(x,r) \in \mathbb{R}^n \times \mathbb{R} \mid g(x) \leq r\}.$$
(2.8)

If *g* is sublinear (i.e., convex and positively homogeneous of degree one), then $\partial_{\varepsilon}g(0) = \partial g(0)$, for all $\varepsilon \ge 0$. If $\tilde{g}(x) = g(x) - k$, $x \in \mathbb{R}^n$, $k \in \mathbb{R}$, then $epi\tilde{g}^* = epig^* + (0, k)$. It is worth nothing that if *g* is sublinear, then

$$epig^* = \partial g(0) \times \mathbb{R}_+.$$
(2.9)

Moreover, if *g* is sublinear and if $\tilde{g}(x) = g(x) - k$, $x \in \mathbb{R}^n$, and $k \in \mathbb{R}$, then

$$epi\tilde{g}^* = \partial g(0) \times [k, \infty). \tag{2.10}$$

Definition 2.5. Let *C* be a closed convex set in \mathbb{R}^n and $x \in C$.

- (1) Let $N_C(x) = \{v \in \mathbb{R}^n \mid \langle v, y x \rangle \leq 0$, for all $y \in C\}$. Then $N_C(x)$ is called the normal cone to *C* at *x*.
- (2) Let $e \ge 0$. Let $N_C^e(x) = \{v \in \mathbb{R}^n \mid \langle v, y x \rangle \le e$, for all $y \in C\}$. Then $N_C^e(x)$ is called the *e*-normal set to *C* at *x*.
- (3) When *C* is a closed convex cone in \mathbb{R}^n , $N_C(0)$ we denoted by C^* and called the negative dual cone of *C*.

Proposition 2.6 (see [17, 18]). Let $f : \mathbb{R}^n \to \mathbb{R}$ be a convex function and let δ_C be the indicator function with respect to a closed convex subset C of \mathbb{R}^n , that is, $\delta_C(x) = 0$ if $x \in C$, and $\delta_C(x) = +\infty$ if $x \notin C$. Let $e \ge 0$. Then

$$\partial_{\varepsilon} (f + \delta_{C})(\overline{x}) = \bigcup_{\substack{e_{0} \ge 0, e_{1} \ge 0\\ e_{0} + e_{1} = \varepsilon}} \{\partial_{e_{0}} f(\overline{x}) + \partial_{e_{1}} \delta_{C}(\overline{x})\}.$$
(2.11)

Proposition 2.7 (see [7]). Let $g : \mathbb{R}^n \to \mathbb{R}$ be a continuous convex function and let $h : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$ be a proper lower semicontinuous convex function. Then

$$epi(g+h)^* = epig^* + epih^*.$$
 (2.12)

Following the proof of Lemma 2.2 in [1], we can prove the following lemma.

Lemma 2.8. Let $F_i \in S_n$, i = 0, 1, ..., m. Suppose that $A \neq \emptyset$. Let $u \in \mathbb{R}^m$ and $\alpha \in \mathbb{R}$. Then the following are equivalent:

3. *e***-Duality Theorem**

Now we give *e*-duality theorems for SDP. Using Lemma 2.8, we can obtain the following lemma which is useful in proving our *e*-strong duality theorems for SDP.

Lemma 3.1. Let $\overline{x} \in A$. Suppose that

$$\bigcup_{(Z,\delta)\in S\times\mathbb{R}_+} \left\{ \begin{pmatrix} \widehat{F}^*(Z) \\ -\operatorname{Tr}[ZF_0] - \delta \end{pmatrix} \right\} - C^* \times \mathbb{R}_+$$
(3.1)

is closed. Then \overline{x} *is an* ϵ *-approximate solution of SDP if and only if there exists* $Z \in S$ *such that for any* $x \in C$ *,*

$$f(x) - \operatorname{Tr}[ZF(x)] \ge f(\overline{x}) - \epsilon. \tag{3.2}$$

Proof. (\Longrightarrow) Let \overline{x} be an ϵ -approximate solution of SDP. Then $f(x) \ge f(\overline{x}) - \epsilon$, for any $x \in A$. Let $h(x) = f(x) - f(\overline{x}) + \epsilon$. Then $h(x) + \delta_A(x) \ge 0$, for any $x \in \mathbb{R}^n$. Thus we have, from Proposition 2.7,

$$0 \in epi(h + \delta_A)^* = epih^* + epi\delta_A^*$$

= $epif^* + (0, f(\overline{x}) - \epsilon) + epi\delta_A^*,$ (3.3)

and hence, $(0, e-f(\overline{x})) \in epif^* + epi\delta_A^*$. So there exists $(u, r) \in epif^*$ such that $(-u, e-f(\overline{x})-r) \in epi\delta_A^*$ and hence there exists $(u, r) \in epif^*$ such that $\langle -u, x \rangle \leq e - f(\overline{x}) - r$ for any $x \in A$. Since $f^*(u) \leq r, \langle -u, x \rangle \leq e - f(\overline{x}) - f^*(u)$ for any $x \in A$; and hence it follows from Lemma 2.8 that

$$\binom{u}{-\epsilon + f(\overline{x}) + f^{*}(u)} \in \bigcup_{(Z,\delta)\in S\times\mathbb{R}_{+}} \left\{ \binom{\widehat{F}^{*}(Z)}{-\operatorname{Tr}[ZF_{0}] - \delta} \right\} - C^{*}\times\mathbb{R}_{+}.$$
(3.4)

Thus there exist $(Z, \delta) \in S \times \mathbb{R}_+, c^* \in C^*$, and $\gamma \in \mathbb{R}_+$ such that

$$u = \widehat{F}^*(Z) - c^*,$$

$$-\epsilon + f(\overline{x}) + f^*(u) = -\operatorname{Tr}[ZF_0] - \delta - \gamma.$$
(3.5)

This gives

$$\left\langle \widehat{F}^{*}(Z), x \right\rangle - \left\langle c^{*}, x \right\rangle - f(x) = \left\langle u, x \right\rangle - f(x) \leq f^{*}(u)$$

$$= -\operatorname{Tr}[ZF_{0}] - \delta - \gamma - f(\overline{x}) + \epsilon,$$
(3.6)

for any $x \in \mathbb{R}^n$. Thus we have

$$f(\overline{x}) - \epsilon \leq -\langle u, x \rangle + f(x) - \operatorname{Tr}[ZF_0] - \delta - \gamma$$

$$= f(x) - \left\langle \widehat{F}^*(Z), x \right\rangle + \langle c^*, x \rangle - \operatorname{Tr}[ZF_0] - \delta - \gamma$$

$$= f(x) - \operatorname{Tr}[ZF(x)] + \langle c^*, x \rangle - \delta - \gamma$$

$$\leq f(x) - \operatorname{Tr}[ZF(x)]$$
(3.7)

for any $x \in C$.

(\Leftarrow) Suppose that there exists $Z \in S$ such that

$$f(x) - \operatorname{Tr}[ZF(x)] \ge f(\overline{x}) - \epsilon, \tag{3.8}$$

for any $x \in C$. Then we have

$$f(x) \ge f(x) - \operatorname{Tr}[ZF(x)] \ge f(\overline{x}) - \epsilon, \tag{3.9}$$

for any $x \in A$. Thus $f(x) \ge f(\overline{x}) - \epsilon$, for any $x \in A$. Hence \overline{x} is an ϵ -approximate solution of SDP.

Now we formulate the dual problem SDD of SDP as follows:

(SDD) maximize
$$f(x) - \operatorname{Tr}[ZF(x)],$$

subject to $0 \in \partial_{\varepsilon_0} f(x) - \hat{F}^*(Z) + N_C^{\varepsilon_1}(x),$
 $Z \succeq 0,$
 $\varepsilon_0 + \varepsilon_1 \in [0, \varepsilon].$
(3.10)

We prove *e*-weak and *e*-strong duality theorems which hold between SDP and SDD.

Theorem 3.2 (ϵ -weak duality). For any feasible solution x of SDP and any feasible solution (y, Z) of SDD,

$$f(x) \ge f(y) - \operatorname{Tr}[ZF(y)] - \epsilon.$$
(3.11)

Proof. Let *x* and (y, Z) be feasible solutions of SDP and SDD respectively. Then $\text{Tr}[ZF(x)] \ge 0$ and there exist $v \in \partial_{e_0} f(y)$ and $\omega \in N_C^{e_1}(y)$ such that $v = -\omega + \hat{F}^*(Z)$. Thus, we have

$$f(x) - \{f(y) - \operatorname{Tr}[ZF(y)]\} \ge \langle v, x - y \rangle - e_0 + \operatorname{Tr}[ZF(y)]$$

$$= \langle -\omega + \widehat{F}^*(Z), x - y \rangle - e_0 + \operatorname{Tr}[ZF(y)]$$

$$\ge \langle \widehat{F}^*(Z), x - y \rangle - e_0 - e_1 + \operatorname{Tr}[ZF(y)]$$

$$= \langle \widehat{F}^*(Z), x \rangle - \langle \widehat{F}^*(Z), y \rangle - e_0 - e_1$$

$$+ \operatorname{Tr}[ZF(y)]$$

$$= \operatorname{Tr}\left[Z\sum_{i=1}^m x_i F_i\right] - \operatorname{Tr}\left[Z\sum_{i=1}^m y_i F_i\right] - e_0 - e_1$$

$$+ \operatorname{Tr}[ZF_0] + \operatorname{Tr}\left[Z\sum_{i=1}^m y_i F_i\right]$$

$$= \operatorname{Tr}[ZF(x)] - e_0 - e_1$$

$$\ge -e_0 - e_1$$

$$\ge -e_0.$$
(3.12)

Hence $f(x) \ge f(y) - \operatorname{Tr}[ZF(y)] - \epsilon$.

Theorem 3.3 (*e*-strong duality). Suppose that

$$\bigcup_{(Z,\delta)\in S\times\mathbb{R}_{+}} \left\{ \begin{pmatrix} \widehat{F}^{*}(Z) \\ -\operatorname{Tr}[ZF_{0}] - \delta \end{pmatrix} \right\} - C^{*}\times\mathbb{R}_{+}$$
(3.13)

is closed. If \overline{x} *is an* ϵ *-approximate solution of SDP, then there exists* $\overline{Z} \in S$ *such that* $(\overline{x}, \overline{Z})$ *is a* 2ϵ *-approximate solution of SDD.*

Proof. Let $\overline{x} \in A$ be an ϵ -approximate solution of SDP. Then $f(x) \ge f(\overline{x}) - \epsilon$, for any $x \in A$. By Lemma 3.1, there exists $\overline{Z} \in S$ such that

$$f(x) - \operatorname{Tr}\left[\overline{Z}F(x)\right] \ge f(\overline{x}) - \epsilon,$$
 (3.14)

for any $x \in C$. Letting $x = \overline{x}$ in (3.14), $\operatorname{Tr}[\overline{Z}F(\overline{x})] \leq e$. Since $F(\overline{x}) \in S$ and $\overline{Z} \in S$, $\operatorname{Tr}[\overline{Z}F(\overline{x})] \geq 0$.

Thus from (3.14),

$$f(x) - \operatorname{Tr}\left[\overline{Z}F(x)\right] + \epsilon \ge f(\overline{x}) - \operatorname{Tr}\left[\overline{Z}F(\overline{x})\right]$$
(3.15)

for any $x \in C$. Hence \overline{x} is an *e*-approximate solution of the following problem:

maximize
$$f(x) - \operatorname{Tr}\left[\overline{Z}F(x)\right]$$
, (3.16)
subject to $x \in C$,

and so, $0 \in \partial_{\epsilon}(f - \hat{F}^*(\overline{Z}) + \delta_C)$ (\overline{x}), and hence, by Proposition 2.6, there exist $\epsilon_0, \epsilon_1 \in [0, \epsilon]$ such that $\epsilon_0 + \epsilon_1 = \epsilon$ and

$$0 \in \partial_{\epsilon_0} f(\overline{x}) - \widehat{F}^* \left(\overline{Z} \right) + N_C^{\epsilon_1}(\overline{x}).$$
(3.17)

So, $(\overline{x}, \overline{Z})$ is a feasible solution of SDD. For any feasible solution (y, Z) of SDD,

$$f(\overline{x}) - \operatorname{Tr}\left[\overline{Z}F(\overline{x})\right] - \left\{f(y) - \operatorname{Tr}\left[ZF(y)\right]\right\} = f(\overline{x}) - \left\{f(y) - \operatorname{Tr}\left[ZF(y)\right]\right\}$$
$$- \operatorname{Tr}\left[\overline{Z}F(\overline{x})\right]$$
$$\geqq -\epsilon - \operatorname{Tr}\left[\overline{Z}F(\overline{x})\right] \qquad (3.18)$$
$$(by \ e\text{-weak duality})$$
$$\geqq -\epsilon - \epsilon$$
$$= -2\epsilon.$$

Thus $(\overline{x}, \overline{Z})$ is a 2 ϵ -approximate solution to SDD.

Now we characterize the ϵ -normal set to \mathbb{R}^{n}_{+} .

Proposition 3.4. Let $(x_1, \ldots, x_n) \in \mathbb{R}^n_+$ and $\epsilon \ge 0$. Then

$$N_{\mathbb{R}^n_+}^{\epsilon}(x_1,\ldots,x_n) = \bigcup_{\substack{\epsilon_i \ge 0\\ \sum_{i=1}^n \epsilon_i = \epsilon}} \prod_{i=1}^n A(\epsilon_i),$$
(3.19)

where

$$A(\epsilon_i) = \begin{cases} -\mathbb{R}_+ & \text{if } x_i = 0, \\ \left[-\frac{\epsilon_i}{x_i}, 0 \right] & \text{if } x_i > 0. \end{cases}$$
(3.20)

Proof. Let $(x_1, \ldots, x_n) \in \mathbb{R}^n_+$ and $\epsilon \ge 0$. Then

$$N_{\mathbb{R}^{n}_{+}}^{e}(x_{1},\ldots,x_{n}) = \partial_{e}\delta_{\mathbb{R}^{n}_{+}}(x_{1},\ldots,x_{n})$$

$$= \partial_{e}\left(\sum_{i=1}^{n}\delta_{\mathbb{R}\times\cdots\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}}\right)(x_{1},\ldots,x_{n})$$

$$= \bigcup_{\substack{e_{i}\geq 0\\ \sum_{i=1}^{n}e_{i}=e}}\sum_{i=1}^{n}\partial_{e_{i}}\delta_{\mathbb{R}\times\cdots\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}\times\mathbb{R}}(x_{1},\ldots,x_{n}).$$
(3.21)

Let $(v_1, \ldots, v_n) \in \partial_{\epsilon_i} \delta_{\mathbb{R} \times \cdots \times \mathbb{R} \times \mathbb{R}_+ \times \mathbb{R} \times \cdots \times \mathbb{R}} (x_1, \ldots, x_n)$ (where \mathbb{R}_+ is at the *i*th position in $\mathbb{R} \times \cdots \times \mathbb{R} \times \mathbb{R}_+ \times \mathbb{R} \times \cdots \times \mathbb{R}$)

$$\Leftrightarrow \text{ for any } (y_1, \dots, y_n) \in \mathbb{R} \times \dots \times \mathbb{R} \times \mathbb{R}_+ \times \mathbb{R} \times \dots \times \mathbb{R}, \\ e_i \geq v_1(y_1 - x_1) + \dots + v_i(y_i - x_i) + \dots + v_n(y_n - x_n), \\ \Leftrightarrow \text{ for any } y_i \in \mathbb{R}_+, \quad e_i \geq v_i(y_i - x_i), \quad v_j = 0, \\ \text{ for } j \in \{1, \dots, n\} \setminus \{i\}, \\ \text{ for } j \in \{1, \dots, n\} \setminus \{i\}, \\ \Leftrightarrow v_i \in \begin{cases} -\mathbb{R}_+, & \text{ if } x_i = 0, \\ \left[-\frac{e_i}{x_i}, 0\right], & \text{ if } x_i > 0, \end{cases} \\ v_j = 0 \text{ for } j \in \{1, \dots, n\} \setminus \{i\}. \end{cases}$$

$$(3.22)$$

Thus, we have

$$N_{\mathbb{R}^{n}_{+}}^{\varepsilon_{1}}(x_{1},\ldots,x_{n}) = \bigcup_{\substack{\epsilon_{i} \geq 0 \\ \sum_{i=1}^{n} \varepsilon_{i}=\epsilon}} \sum_{i=1}^{n} \{0\} \times \cdots \times \{0\} \times A(\varepsilon_{i}) \times \{0\} \times \cdots \times \{0\}$$

$$= \bigcup_{\substack{\epsilon_{i} \geq 0 \\ \sum_{i=1}^{n} \varepsilon_{i}=\epsilon}} \prod_{i=1}^{n} A(\varepsilon_{i}).$$
(3.23)

From Proposition 3.4, we can calculate $N^{\epsilon}_{\mathbb{R}^2_+}.$

Corollary 3.5. Let $(x_1, x_2) \in \mathbb{R}^2_+$ and $\epsilon \ge 0$. Then following hold.

(i) If (x₁, x₂) = (0,0), then N^e_{ℝ²+}(0,0) = -ℝ²₊.
(ii) If (x₁, x₂) = (x₁,0) and x₁ > 0, then N^e_{ℝ²+}(x₁,0) = [-ε/x₁,0] × (-∞,0].
(iii) If (x₁, x₂) = (0, x₂) and x₂ > 0, then N^e_{ℝ²+}(0, x₂) = (-∞,0] × [-ε/x₂,0].
(iv) If (x₁, x₂) = (x₁, x₂), and x₁ > 0 and x₂ > 0, then

$$N_{\mathbb{R}^{2}_{+}}^{e}(x_{1}, x_{2}) = \bigcup_{\substack{e_{1} \ge 0, e_{2} \ge 0, \\ e_{1} + e_{2} = e}} \left[-\frac{e_{1}}{x_{1}}, 0 \right] \times \left[-\frac{e_{2}}{x_{2}}, 0 \right].$$
(3.24)

Now we give an example illustrating our *e*-duality theorems. *Example 3.6.* Consider the following convex semidefinite program.

(SDP) Minimize
$$x_1 + x_2^2$$
,
subject to $\begin{pmatrix} 0 & x_1 \\ x_1 & 0 \end{pmatrix} \ge 0$, (3.25)
 $(x_1, x_2) \in \mathbb{R}^2_+$.

Let $f(x_1, x_2) = x_1 + x_2^2$,

$$F_{0} = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \qquad F_{1} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \qquad F_{2} = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \qquad (3.26)$$

and $\epsilon \ge 0$. Let $f(x_1, x_2) = x_1 + x_2^2$ and

$$F(x_1, x_2) = \begin{pmatrix} 0 & x_1 \\ x_1 & 0 \end{pmatrix}.$$
 (3.27)

Then $A := \{(0, x_2) \in \mathbb{R}^2 \mid x_2 \ge 0\}$ is the set of all feasible solutions of SDP and the set of all ϵ -approximate solutions of SDP is $\{(x_1, x_2) \in \mathbb{R}^2 \mid x_1 = 0, 0 \le x_2 \le \sqrt{\epsilon}\}$. Let $F = \{((x_1, x_2), Z) \mid 0 \in \partial_{\epsilon_0} f(x_1, x_2) - \widehat{F}^*(Z) + N_{\mathbb{R}^2_+}^{\epsilon_1}(x_1, x_2), Z \ge 0, \epsilon_0 + \epsilon_1 \in [0, \epsilon]\}$. Then F is the set of all feasible solution of SDD. Now we calculate the set F.

$$\begin{split} \tilde{A} &:= \left\{ \left((0,0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid 0 \in \partial_{e_0} f(0,0) - \tilde{F}^* \begin{pmatrix} a & b \\ b & c \end{pmatrix} + N_{\mathbb{R}^2_2}^{e_1}(0,0), \\ a &\geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid 0 \in [1] \times [-2\sqrt{e_0}, 2\sqrt{e_0}] - (2b,0) - \mathbb{R}^2_+, \\ a &\geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid (2b,0) \in (-\infty,1] \times (-\infty,2\sqrt{e_0}], \\ a &\geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid a \geq 0, \ c \geq 0, \ b \leq \frac{1}{2}, \ b^2 \leq ac \right\}, \\ \tilde{B} &:= \left\{ \left((0,x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_2 > 0, \ 0 \in \partial_{e_0} f(0,x_2) - \tilde{F}^* \begin{pmatrix} a & b \\ b & c \end{pmatrix} + N_{\mathbb{R}^2_+}^{e_0}(0,x_2) \\ &= 20, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_2 > 0, \ 0 \in \partial_{e_0} f(0,x_2) - \tilde{F}^* \begin{pmatrix} a & b \\ b & c \end{pmatrix} + N_{\mathbb{R}^2_+}^{e_0}(0,x_2) \\ &= 20, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_2 > 0, \ 0 \in [1] \times [2x_2 - 2\sqrt{e_0}, 2x_2 + 2\sqrt{e_0}] - (2b,0) \\ &+ (-\infty,0] \times \left[-\frac{e_1}{x_2}, 0 \right], \ a \geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_2 > 0, \ (2b,0) \in (-\infty,1] \times \left[2x_2 - \frac{e_1}{x_2} - 2\sqrt{e_0}, 2x_2 + 2\sqrt{e_0} \right], \\ &= 20, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_2 > 0, \ (2b,0) \in (-\infty,1] \times \left[2x_2 - \frac{e_1}{x_2} - 2\sqrt{e_0}, 2x_2 + 2\sqrt{e_0} \right], \\ &= 20, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0,e] \right\} \\ &= \left\{ \left((0,x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid 0 < x_2 \leq \frac{\sqrt{e_0} + \sqrt{e_0 + 2e_1}}{2}, \ a \geq 0, \ c \geq 0, \ b \leq \frac{1}{2}, \ b^2 \leq ac, \\ e_0 + e_1 \in [0,e] \right\}, \end{split} \right\}$$

$$\begin{split} \tilde{C} &:= \left\{ \left((x_1, 0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_1 > 0, \ 0 \in \partial_{e_0} f(x_1, 0) - \hat{F}^* \begin{pmatrix} a & b \\ b & c \end{pmatrix} + N_{\mathbb{R}^2_+}^{e_1}(x_1, 0), \\ a &\geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0, e] \right\} \\ &= \left\{ \left((x_1, 0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_1 > 0, \ 0 \in \{1\} \times [-2\sqrt{e_0}, +2\sqrt{e_0}] - (2b, 0) \\ &+ \left[-\frac{e_1}{x_1}, 0 \right] \times (-\infty, 0], \ a \geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0, e] \right\} \\ &= \left\{ \left((x_1, 0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_1 > 0, \ (2b, 0) \in \left[1 - \frac{e_1}{x_1}, 1 \right] \times (-\infty, 2\sqrt{e_0}], \\ a \geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0, e] \right\} \\ &= \left\{ \left((x_1, 0), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid 0 < x_1, \ -e_1 \leq -x_1 + 2bx_1, \ a \geq 0, \ c \geq 0, \ b \leq \frac{1}{2}, \ b^2 \leq ac, \\ e_0 + e_1 \in [0, e] \right\}, \\ \tilde{D} := \left\{ \left((x_1, x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid 0 < x_1, \ -e_1 \leq -x_1 + 2bx_1, \ a \geq 0, \ c \geq 0, \ b \leq \frac{1}{2}, \ b^2 \leq ac, \\ e_0 + e_1 \in [0, e] \right\}, \\ \tilde{D} := \left\{ \left((x_1, x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_1 > 0, \ x_2 > 0, \ 0 \in \partial_{e_0} f(x_1, x_2) - \hat{F}^* \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \\ &+ N_{e_1^2}^{e_1}(x_1, x_2), \ a \geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_0 + e_1 \in [0, e] \right\} \\ &= \left\{ \left((x_1, x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_1 > 0, \ x_2 > 0, \ 0 \in \{1\} \times [2x_2 - 2\sqrt{e_0}, 2x_2 + 2\sqrt{e_0}] - (2b, 0) + \left[-\frac{e_1^1}{x_1}, 0 \right] \times \left[-\frac{e_1^2}{x_2}, 0 \right], \\ a \geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_1^1 + e_1^2 = e_1, \ e_0 + e_1 \in [0, e] \right\} \\ &= \left\{ \left((x_1, x_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid x_1 > 0, \ x_2 > 0, \\ (2b, 0) \in \left[1 - \frac{e_1^1}{x_1}, 1 \right] \times \left[2x_2 - \frac{e_1^2}{x_2} - 2\sqrt{e_0}, 2x_2 + 2\sqrt{e_0} \right], \\ a \geq 0, \ c \geq 0, \ b^2 \leq ac, \ e_1^1 + e_1^2 = e_1, \ e_0 + e_1 \in [0, e] \right\} \end{split}$$

$$= \left\{ \left((x_{1}, x_{2}), \begin{pmatrix} a & b \\ b & c \end{pmatrix} \right) \mid 0 < x_{1}, -\epsilon_{1}^{1} \leq -x_{1} + 2bx_{1}, \\ 0 < x_{2} \leq \frac{\sqrt{\epsilon_{0}} + \sqrt{\epsilon_{0} + 2\epsilon_{1}^{2}}}{2}, \ a \geq 0, \ c \geq 0, \ b \leq \frac{1}{2}, \ b^{2} \leq ac, \ \epsilon_{1}^{1} + \epsilon_{1}^{2} = \epsilon_{1}, \\ \epsilon_{0} + \epsilon_{1} \in [0, \epsilon] \right\}.$$
(3.28)

Thus $F = \tilde{A} \cup \tilde{B} \cup \tilde{C} \cup \tilde{D}$. We can check that for any $(x_1, x_2) \in A$ and any $((y_1, y_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix}) \in F$,

$$f(x_1, x_2) \ge f(y_1, y_2) - \operatorname{Tr}\left(\begin{pmatrix} a & b \\ b & c \end{pmatrix} F(y_1, y_2) \right) - \epsilon,$$
(3.29)

that is, *e*-weak duality holds.

Let $(\overline{x}_1, \overline{x}_2) \in A$ be an *e*-approximate solution of SDP. Then $\overline{x}_1 = 0$ and $0 \leq \overline{x}_2 \leq \sqrt{\epsilon}$. So, we can easily check that $((\overline{x}_1, \overline{x}_2), \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}) \in F$. Since $\operatorname{Tr}(\begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix} F(\overline{x}_1, \overline{x}_2)) = 0$, from (3.29),

$$f(\overline{x}_1, \overline{x}_2) \ge f(y_1, y_2) - \operatorname{Tr}\left(\begin{pmatrix} a & b \\ b & c \end{pmatrix} F(y_1, y_2) \right) - \epsilon,$$
(3.30)

for any $((y_1, y_2), \begin{pmatrix} a & b \\ b & c \end{pmatrix}) \in F$. So $((\overline{x}_1, \overline{x}_2), \begin{pmatrix} \overline{a} & \overline{b} \\ \overline{b} & \overline{c} \end{pmatrix})$ is an ϵ -approximate solution of SDD. Hence ϵ -strong duality holds.

Acknowledgment

This work was supported by the Korea Science and Engineering Foundation (KOSEF) NRL Program grant funded by the Korean government (MEST) (no. R0A-2008-000-20010-0).

References

- [1] V. Jeyakumar and N. Dinh, "Avoiding duality gaps in convex semidefinite programming without Slater's condition," Applied Mathematics Report AMR04/6, University of New South Wales, Sydney, Australia, 2004.
- [2] M. V. Ramana, L. Tuncel, and H. Wolkowicz, "Strong duality for semidefinite programming," SIAM Journal on Optimization, vol. 7, no. 3, pp. 641-662, 1997.
- [3] V. Jeyakumar, G. M. Lee, and N. Dinh, "New sequential Lagrange multiplier conditions characterizing optimality without constraint qualification for convex programs," SIAM Journal on Optimization, vol. 14, no. 2, pp. 534–547, 2003.
- [4] V. Jeyakumar and M. J. Nealon, "Complete dual characterizations of optimality for convex semidefinite programming," in Constructive, Experimental, and Nonlinear Analysis (Limoges, 1999), vol. 27 of CMS Conference Proceedings, pp. 165–173, American Mathematical Society, Providence, RI, USA, 2000.

- [5] N. Dinh, V. Jeyakumar, and G. M. Lee, "Sequential Lagrangian conditions for convex programs with applications to semidefinite programming," *Journal of Optimization Theory and Applications*, vol. 125, no. 1, pp. 85–112, 2005.
- [6] V. Jeyakumar, G. M. Lee, and N. Dinh, "Lagrange multiplier conditions characterizing the optimal solution sets of cone-constrained convex programs," *Journal of Optimization Theory and Applications*, vol. 123, no. 1, pp. 83–103, 2004.
- [7] V. Jeyakumar, G. M. Lee, and N. Dinh, "Characterizations of solution sets of convex vector minimization problems," *European Journal of Operational Research*, vol. 174, no. 3, pp. 1380–1395, 2006.
- [8] M. G. Govil and A. Mehra, "ε-optimality for multiobjective programming on a Banach space," European Journal of Operational Research, vol. 157, no. 1, pp. 106–112, 2004.
- [9] C. Gutiérrez, B. Jiménez, and V. Novo, "Multiplier rules and saddle-point theorems for Helbig's approximate solutions in convex Pareto problems," *Journal of Global Optimization*, vol. 32, no. 3, pp. 367–383, 2005.
- [10] A. Hamel, "An ε-lagrange multiplier rule for a mathematical programming problem on Banach spaces," Optimization, vol. 49, no. 1-2, pp. 137–149, 2001.
- [11] V. Jeyakumar and B. M. Glover, "Characterizing global optimality for DC optimization problems under convex inequality constraints," *Journal of Global Optimization*, vol. 8, no. 2, pp. 171–187, 1996.
- [12] G. S. Kim and G. M. Lee, "On ε-approximate solutions for convex semidefinite optimization problems," *Taiwanese Journal of Mathematics*, vol. 11, no. 3, pp. 765–784, 2007.
- [13] J. C. Liu, "ε-duality theorem of nondifferentiable nonconvex multiobjective programming," Journal of Optimization Theory and Applications, vol. 69, no. 1, pp. 153–167, 1991.
- [14] J. C. Liu, "ε-Pareto optimality for nondifferentiable multiobjective programming via penalty function," Journal of Mathematical Analysis and Applications, vol. 198, no. 1, pp. 248–261, 1996.
- [15] J.-J. Strodiot, V. H. Nguyen, and N. Heukemes, "ε-optimal solutions in nondifferentiable convex programming and some related questions," *Mathematical Programming*, vol. 25, no. 3, pp. 307–328, 1983.
- [16] K. Yokoyama and S. Shiraishi, "An ε -optimal condition for convex programming problems without Slater's constraint qualifications," preprint.
- [17] J. B. Hiriart-Urruty and C. Lemarechal, Convex Analysis and Minimization Algorithms. I. Fundamentals, vol. 305 of Grundlehren der mathematischen Wissenschaften, Springer, Berlin, Germany, 1993.
- [18] J. B. Hiriart-Urruty and C. Lemarechal, Convex Analysis and Minimization Algorithms. II. Advanced Theory and Bundle Methods, vol. 306 of Grundlehren der mathematischen Wissenschaften, Springer, Berlin, Germany, 1993.